

Quantum Key Distribution for Railway

Initial Presentation

Sebastian Staib

Design IT. Create Knowledge.



Motivation



- Quantum Computers (QC) that can break currently used asymmetric crypto
 - are expected to be available in ~ 10 years. [1]
- Shor's algorithm enables prime factorization and the calculation of discrete logarithms (in cyclic groups) in polynomial [2]
 - Breaks the security of several public-key methods used today:
 - RSA [3]
 - Diffie-Hellman (ECDH, DSA/ECDSA) [4]
- Adversaries can record encrypted communications for encryption once QC become available ("Harvest Now, Decrypt Later") [5]
- [1]
- [2]
- [3]

Two Approaches to Quantum-Secure Systems



- Post-Quantum Cryptography (PQC):
 - Based on mathematical problems
 - Resistant to both conventional and quantum cryptanalysis [6]
 - Standards are available since August 2024, by NIST: Kyber, Dilithium and SPHINCS+ [7]



- Quantum Key Distribution (QKD):
 - Two parties establish a shared secret key by encoding information in quantum states [8]
 - Generates information-theoretically secure key material that can be used to generate perfectly secure ciphertexts, e.g. One-Time-Pad [9]
- [6]

QKD Architecture



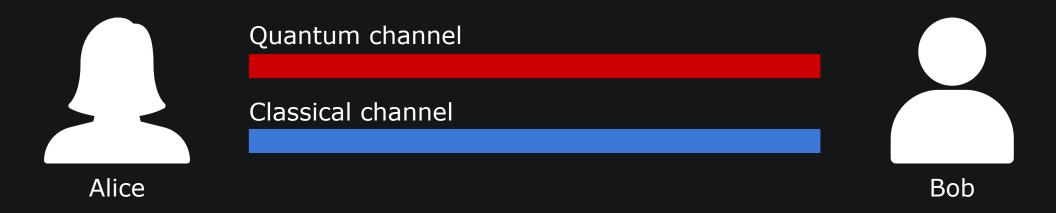
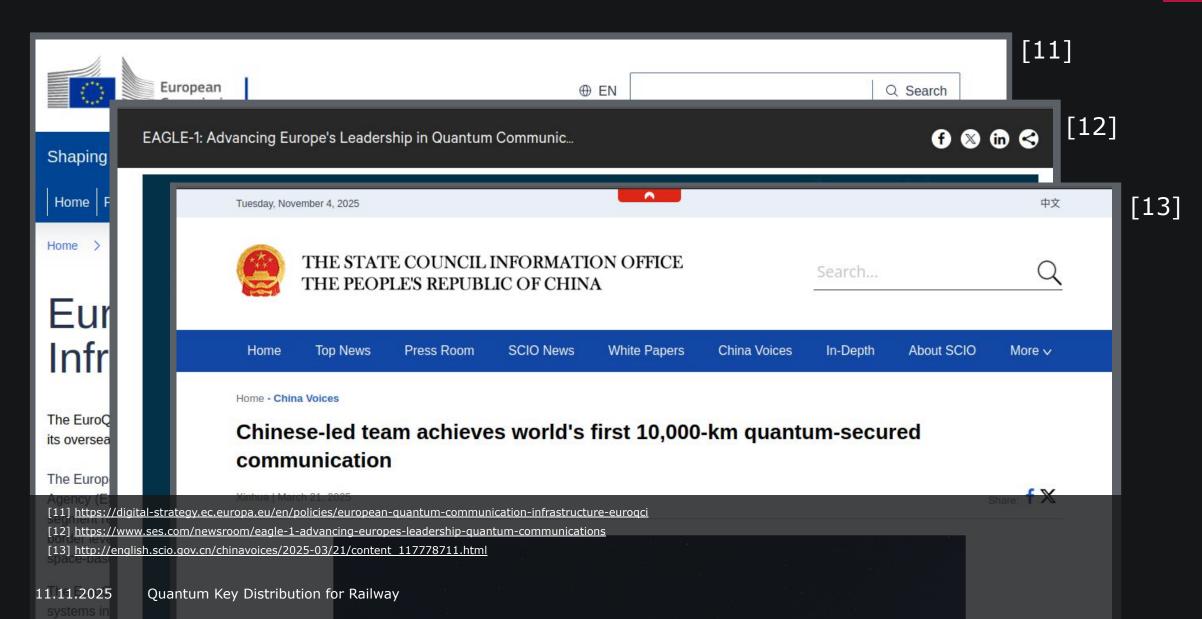


Fig 1. QKD setting. Alice and Bob are connected through a quantum channel and a classical channel [10]

[10] based on Fig 1 in https://doi.org/10.1103/RevModPhys.81.1301

Recent QKD Developments





Content of the Thesis



- Discuss the theoretical and practical foundations of QKD
- Introduce architecture of QKD Networks (QKDN)
- Develop a taxonomy of the various protocols and standards
 - QKD protocols: BB84, E91 and BBM92
 - Different types of QKD (CV-QKD vs. DV-QKD)
 - Standards: *ETSI GS QKD 014* [14] and *ITU-T Y.3800* [15]
- State of the art and related work
- Proof of concept and software architecture
 - Implementation for use case "Schlüsseltankstelle" using the above standards

^[15]

Standards and Reference Architecture[14]

ETSI GS QKD 014 V1.1.1 (2019-02)



Quantum Key Distribution (QKD); Protocol and data format of REST-based key delivery API

[14] https://www.etsi.org/deliver/etsi_gs/QKD/001_099/014/01.01.01_60/gs_qkd014v010101p.pdf

[15] https://www.itu.int/rec/T-REC-Y.3800

[15

International Telecommunication Unior

ITU-T Y.3800

STANDARDIZATION SECTOR
OF ITU

ITU-T Y.3801

ITU-T Y.3802

ITU-T Y.3803

ITU-T Y.3804

TELECOMMUNICATION STANDARDIZATION SECTOR OF ITU (09/2020)

(10/2019)

SERIES Y: GLOBAL INFORMATION
INFRASTRUCTURE, INTERNET PROTOCOL ASPECTS,
NEXT-GENERATION NETWORKS, INTERNET OF
THINGS AND SMART CITIES

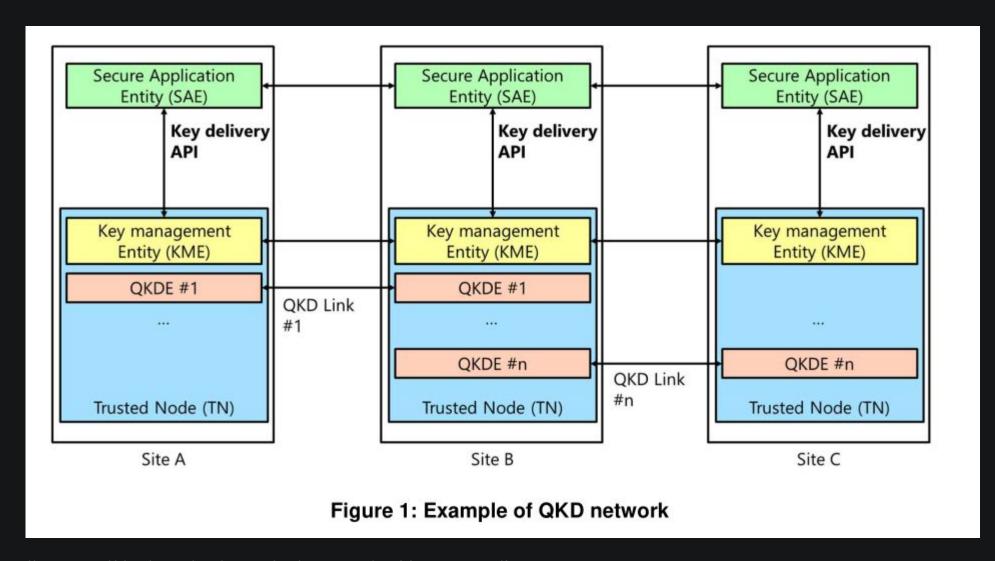
Quantum key distribution networks

Quantum key distribution networks – Control and management

Standards & reference architecture



[14



[14] https://www.etsi.org/deliver/etsi_gs/QKD/001_099/014/01.01.01_60/gs_qkd014v010101p.pdf

Standards & reference architecture



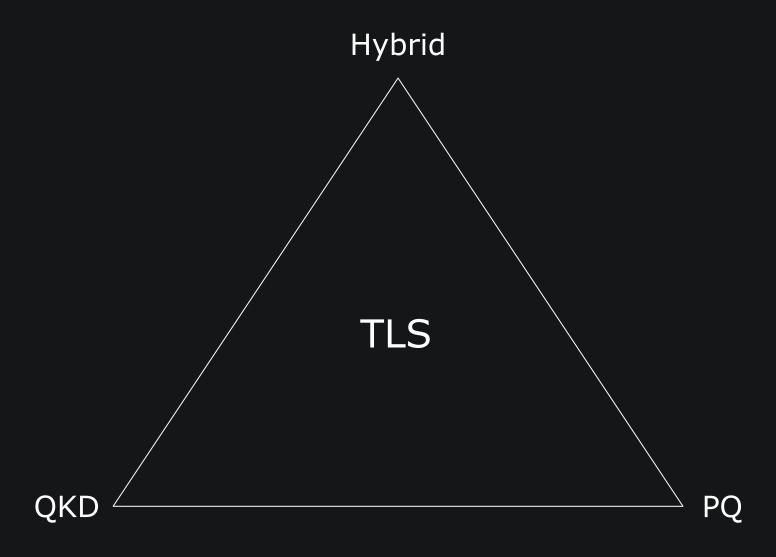
[15]

Cipher text Application sender Application receiver 01001011 Plain text Plain text 11111111 111111111 Encryptor Decryptor Application link Key Key 10110100 10110100 Classical channel QKD module QKD module Quantum channel Y.3800(19)_F01 QKD link Figure 1 – Configuration example of QKD use for securing a P-to-P application link

[15] https://www.itu.int/rec/T-REC-Y.3800

TLS Post-Quantum Approaches (1)

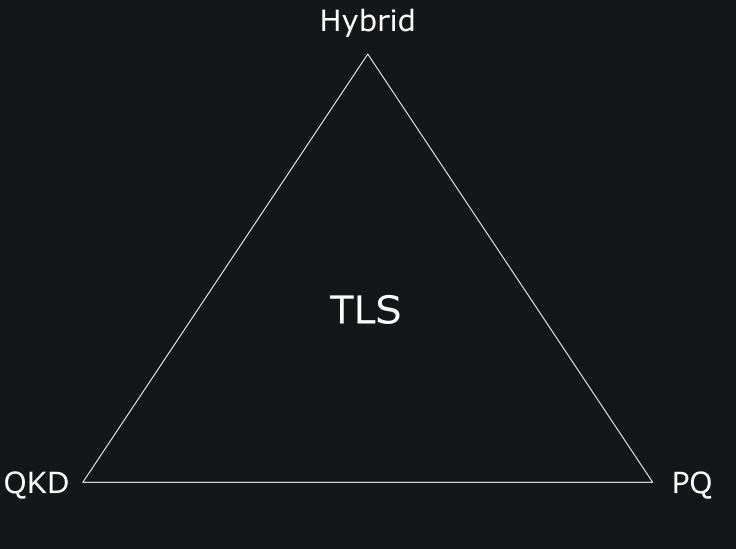




TLS Post-Quantum Approaches (5)



- Evaluation criteria:
 - Security model
 - QC-resistant
 - Handshake/latency behavior
 - mobile contexts
 - Deployability & maturity



TLS Post-Quantum Approaches (6)

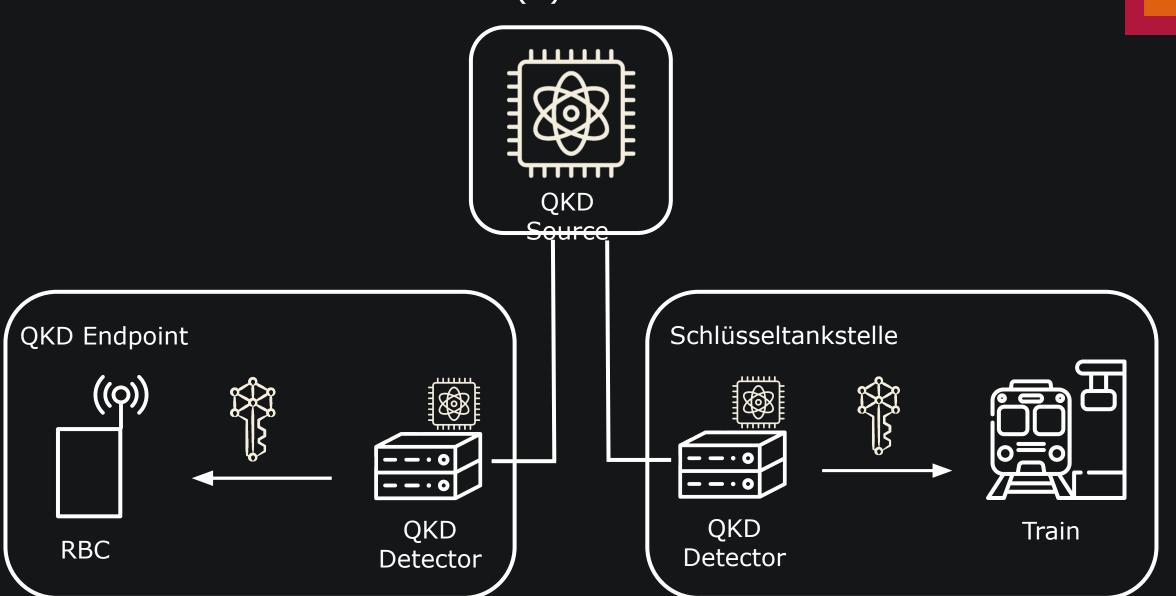


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Table4 V Scheme / Property		TLS-QKD (as described in 2025_Prévost_An ETSI GS QKD compliant TLS)	KEMTLS (as described in 2022_Schwabe_Stebila_Wiggers_Post-quan tum_TLS_without_handshake_signatures)	KEM-QKD (as discussed in 2025_Blanco_Romero_et_alQKD-KEM Hybrid_QKD_Integration_into_TLS_with_Op enSSL_Providers)
Summary	of both the handshake (via transcript bound	Plug quantum keys from ETSI GS QKD 014 (KME/SAE architecture) into ordinary TLS so that endpoints can use QKD-provided entropy without breaking the web-PKI/TLS stack.	uses key-encapsulation mechanisms (KEMs) instead of signatures for server authentication PKI certifies KEM keys; same RTT; smaller/faster than PQ-signed TLS.	Mix a normal post-quantum KEM (e.g., ML-KEM/Kyber) with QKD material so that the TLS 1.3 handshake stays compatible with OpenSSL while the final handshake secret is protected by both PQC and QKD. If either one holds, the session key is safe
Repository		https://github.com/qursa-uc3m/qkd-kem-provider	https://github.com/thomwiggers/kemtls-experiment	QKD-KEM Provider (OpenSSL 3): https://github.co QKD ETSI API (C wrapper): https://github.com/qu QKD-KEM Benchmarking Suite: https://github.cor
Confidentiality	AEAD ciphers at the record layer; the handshake is encrypted from ServerHello onward, so identities and key shares exchanged after that point are confidential. TLS does not hide	feeds QKD-derived symmetric keys into an otherwise standard TLS 1.3 record layer that uses AEAD (e.g., AES-256-GCM), so the channel's confidentiality remains computational (not information-theoretic). AES in AEAD mode with a 256-bit key ("hardcoded" key size in the prototype)	Similar to classical TLS 1.3	Similar to classical TLS 1.3
Availabilty	Replay Attacks: 0-RTT replay, where an attacker duplicates early-data flights or exploits inconsistent server-state/blocked ServerHello to make the same request run multiple times, overloading rate-limiters and backends. Attack works by replaying ClientHello/early data at scale, polluting or bypassing replay caches if misconfigured, and abusing client retry behavior in multi-zone deployments. [4]	Same as classical TLS 1.3	No difference from classical TLS 1.3	No difference from classical TLS 1.3
Integrity	yes [4]	yes, from AEAD just like clasiscal TLS 1.3	Uses AEAD	Uses AEAD
Forward Secrecy	yes [4]	yes [22]	Provided by ephemeral KEM exchange	Provided by ephemeral KEM exchange
		Yes, Possesion of the same QKD key proven with		

Use Case: "Schlüsseltankstelle" (1)

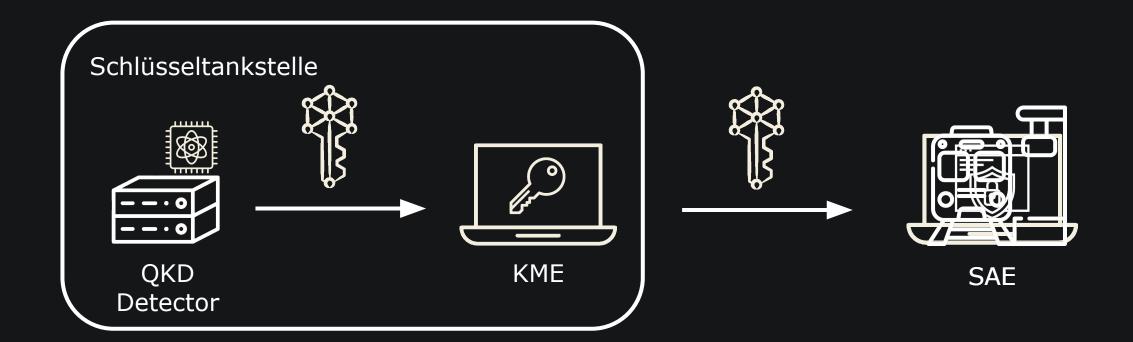




11.11.2025

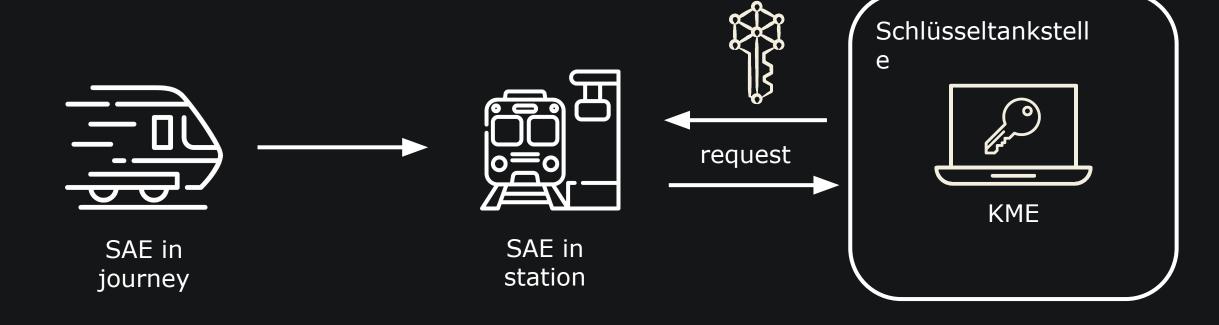
Use Case: "Schlüsseltankstelle" (2)





Use Case: "Schlüsseltankstelle" (3)

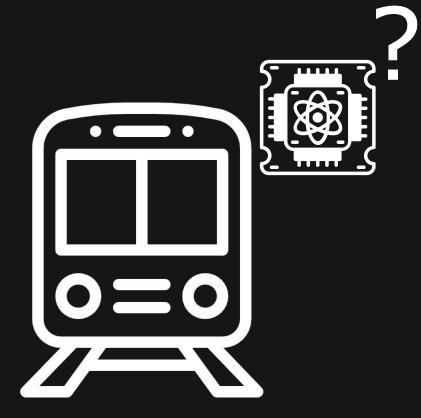




Open Questions







Quantum optics required on the train?



What about key storage capacity?

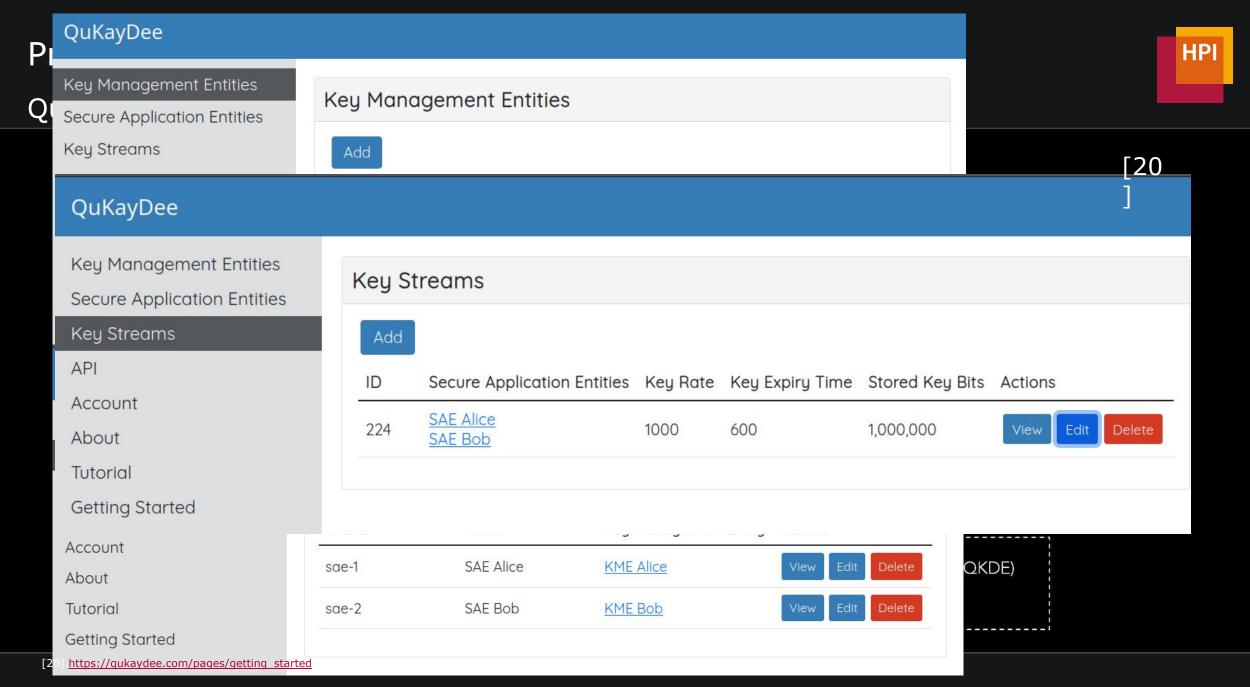
Source: Flaticon.com

Prototype Plan



• Phase 1

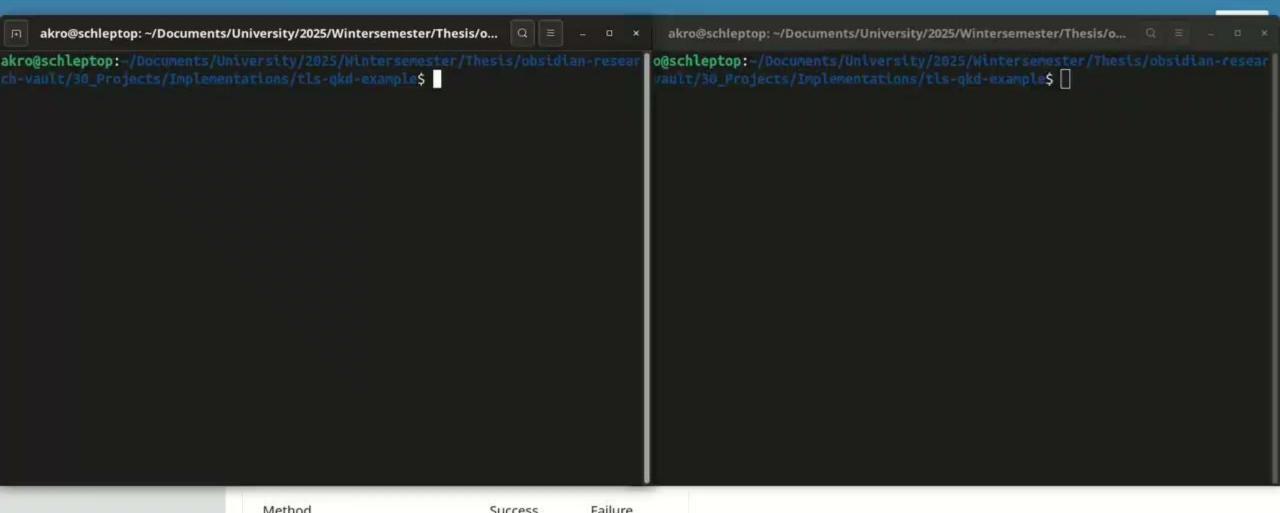
- API-level integration with "Hello World" and test harness using QKD key delivery interfaces (ETSI GS QKD 014 v1.1.1)
- using Key Material from QuKayDee [19]
- Phase 2
 - Evalutiion phase of different TLS variants and decision based on evaluation criteria
- Phase 3
 - Prototype implementation with chosen variant for key management solution, specifically addressed to Mobile Use Case "Schlüsseltankstelle".



Proof of Concept – Approach One (2)



Video



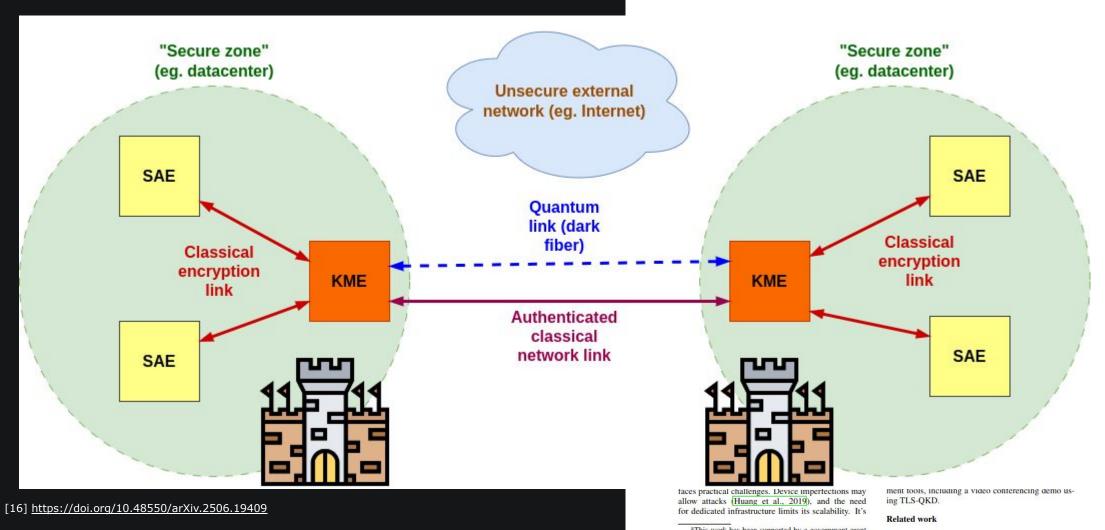
Wictiod	Success	randic
Get status	2	0
Get key	7	0

Proof of Concept – Approach Two (1)



[16]

TLS-QKD Topology



11.11.2025 Quantum Key Distribution for Railway

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Proof of Concept – Approach Two (2)

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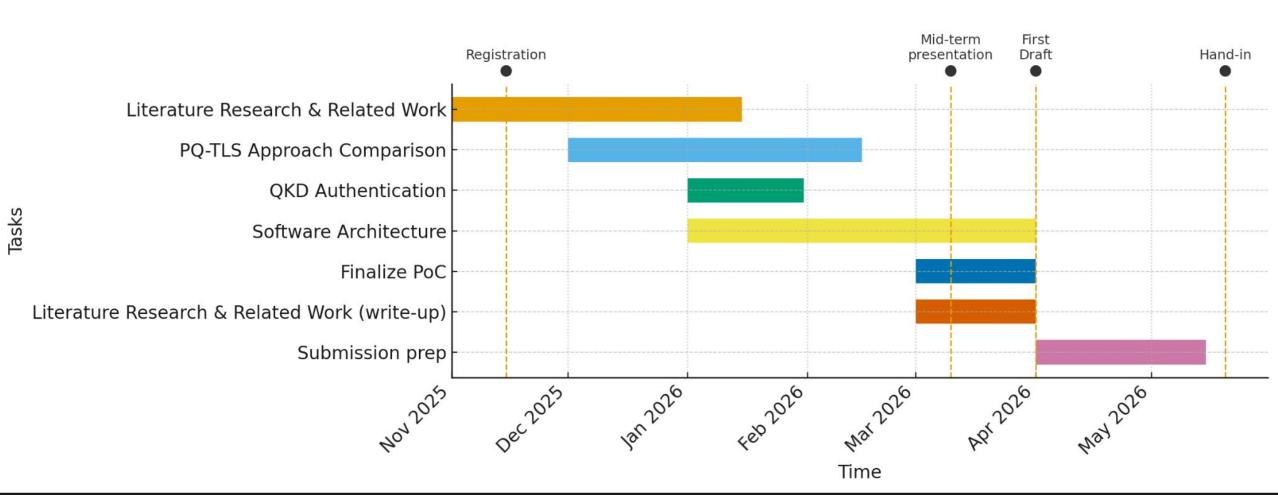
Video



Timetable







Working Packages



- 1. What are different post quantum secure alternatives for TLS? How do they perform compared to other approaches (e.g. KEMTLS)? Compare on security, authentication, latency, deployment model, maturity level, hardware requirements, resistance to QC and the mobile use case "Schlüsseltankstelle".
- 2. Trust Bootstrapping and Authentication: What is the abstract concept of authentication in QKDN? What approaches exist to establish trust between two parties? What are their differences? (PKI vs. PSK)
- 3. Architecture mapping to the "Schlüsseltankstelle" Use Case: How can a high-level QKDN architecture be mapped to this use case? Clarify roles, responsibilities, interfaces at a conceptual level, the policies that govern key use and distribution



Backup Slides

Design IT. Create Knowledge.

Introduction



Master student: Cybersecurity

Work experience:

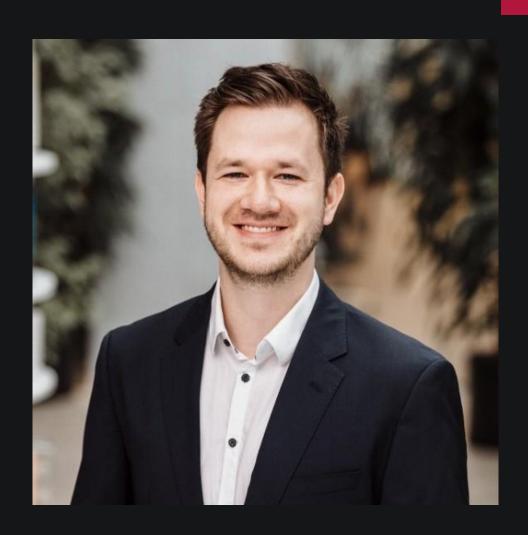
Worked in IT-Consulting

"IT-Grundschutzpraktiker" (certified)

Incident response for ~1 year,Focus on crisis management

Academic Progress:

- Key areas of cybersecurity
 - Network and application security
- Focus on identity management and applied cryptography
- Thesis about Quantum Key Distribution



11.11.2025

TLS Post-Quantum Approaches (2)

HPI

An ETSI GS QKD compliant TLS implementation^a

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Keywords: TLS, Quantum Key Distribution, Rust, ETSI.

Abstract: A modification of the TLS protocol is presented, using our implementation of the Quantum Key Distribution

> (QKD) standard ETSI GS QKD 014 v1.1.1. We rely on the Rustls library for this. The TLS protocol is modified while maintaining backward compatibility on the client and server side. We thus wish to participate in the effort to generalize the use of QKD on the Internet. We used our protocol for a video conference call encrypted by QKD. Finally, we analyze the performance of our protocol, comparing the time needed to

establish a handshake to that of TLS 1.3.

1 INTRODUCTION

Quantum computers threaten current public key cryptosystems like RSA and ECC, which are expected to be broken once such machines are operational Bhatia and Ramkumar, 2020). This has prompted concerns about "harvest now, decrypt later" attacks, where adversaries store encrypted data to decrypt in the future (Paul, 2022).

Post-quantum cryptography offers alternatives based on quantum-resistant problems, but new attacks continue to emerge (Kaluderovic, 2022), raising

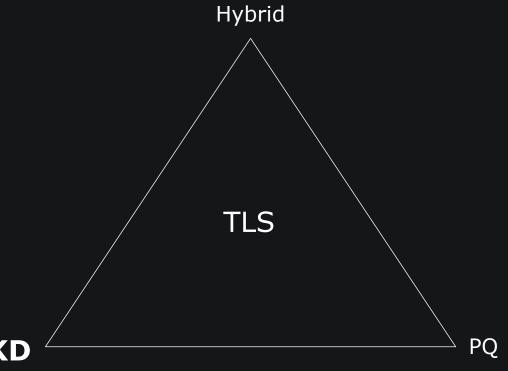
[16] https://doi.org/10.48550/arXiv.2506.19409

best suited for high-security environments like interdatacenter links or governmental networks.

Due to fiber loss and the no-cloning theorem, QKD is limited to a few hundred kilometers (Huttner et al., 2022). Multipath QKD protocols address this with trusted intermediaries (Liu et al., 2024; Prévost et al., 2025). ETSI GS QKD 014 v1.1.1 defines a standard interface for managing QKD keys (ETSI, 2019), which we previously verified with ProVerif under specific assumptions (Prévost et al., 2024).

We present a practical implementation of this standard by integrating QKD into TLS. Our "TLS-Quantum Key Distribution for Railway

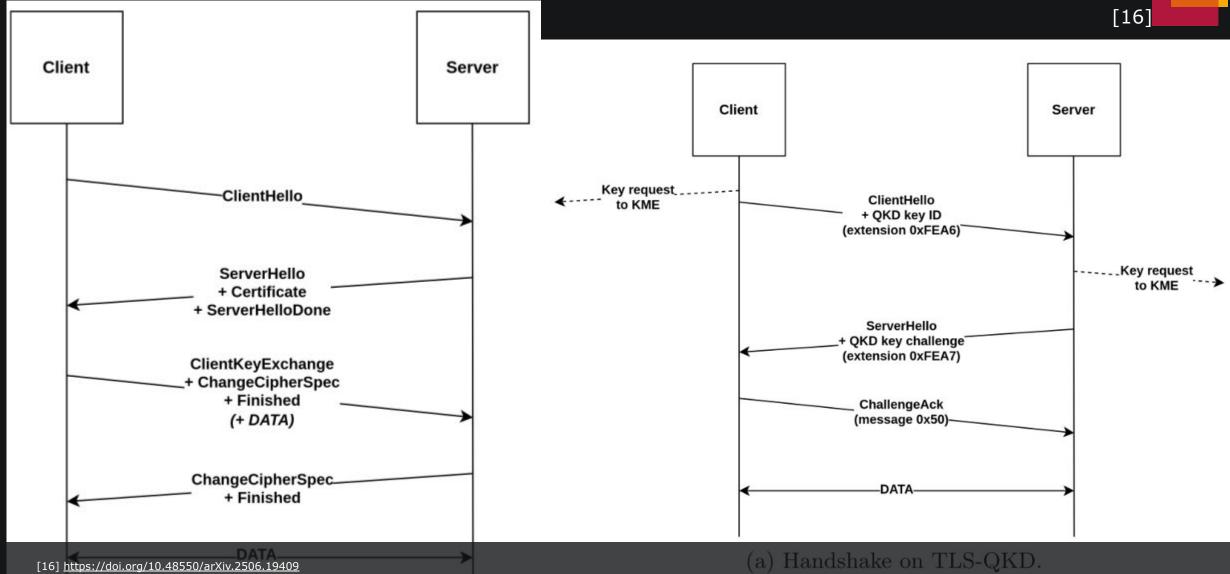
[16]



QKD

"Classical" TLS vs. TLS-QKD





TLS Post-Quantum Approaches (3)

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[17]

Post-Quantum TLS Without Handshake Signatures

Full version, March 15, 2022

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ABSTRACT

We present KEMTLS, an alternative to the TLS 1.3 handshake that uses key-encapsulation mechanisms (KEMs) instead of signatures for server authentication. Among existing post-quantum candidates, signature schemes generally have larger public key/signature sizes compared to the public key/ciphertext sizes of KEMs: by using an IND-CCA-secure KEM for server authentication in post-quantum TLS, we obtain multiple benefits. A size-optimized post-quantum instantiation of KEMTLS requires less than half the bandwidth of a size-optimized post-quantum instantiation of TLS 1.3. In a speed-optimized instantiation, KEMTLS reduces the amount of server CPU cycles by almost 90% compared to TLS 1.3, while at the same time reducing communication size, reducing the time until the client can start sending encrypted application data, and eliminating code for signatures from the server's trusted code base.

Update: in Appendix F we present updated measurements based on the NIST standardization project's round-3 candidate schemes.

CCS CONCEPTS

 Security and privacy → Security protocols; Web protocol security; Public key encryption.

KEYWORDS

Post-quantum cryptography; key-encapsulation mechanisms; Transport Layer Security; NIST PQC

ACM Reference Format:

Peter Schwabe, Douglas Stebila, and Thom Wiggers. 2020. Post-Quantum TLS Without Handshake Signatures: Full version. March 15, 2022. In 2020

[17] https://doi.org/10.1145/3372297.3423350

1 INTRODUCTION

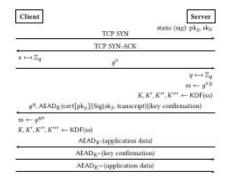


Figure 1: High-level overview of TLS 1.3, using signatures for server authentication.

RSA or elliptic-curve signatures. Public keys for the signatures are embedded in certificates and transmitted during the handshake. Figure 1 gives a high-level overview of the TLS 1.3 protocol, focusing on the signed-Diffie–Hellman aspect of the handshake.

Preparing for post-quantum TLS. There have been many experiments and much research in the past five years on moving the TLS ecosystem to post-quantum cryptography. Most of the work has focused on adding post-quantum key exchange to TLS, usually in the context of so-called "hybrid" key exchange that uses both a post-quantum algorithm and a traditional (usually elliptic curve)

lgorithm, beginning with an experimental demonstration in 201 of ring-LWE-based key exchange in TLS 1.2 [21].

Public experiments by industry started in 2016 with the CECPQ1 experiment by Google [76], combining X25519 ECDH [9] with NewHope lattice-based key exchange [2] in the TLS 1.2 handshake.

A CECPQ2 followup experiment with TLS 1.3 was announced in late Quantum Key Distribution for Railway is currently being run by Google using a combina-

TLS PC

TLS Post-Quantum Approaches (4)

HPI

QKD-KEM: Hybrid QKD Integration into TLS with **OpenSSL Providers**

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[18]

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Abstract-Quantum Key Distribution (QKD) promises information-theoretic security, yet integrating QKD into existing protocols like TLS remains challenging due to its fundamentally different operational model. In this paper, we propose a hybrid QKD-KEM protocol with two distinct integration approaches: a client-initiated flow compatible with both ETSI 004 and 014 specifications, and a server-initiated flow similar to existing work but limited to stateless ETSI 014 APIs. Unlike previous implementations, our work specifically addresses the integration of stateful QKD key exchange protocols (ETSI 004) which is essential for production QKD networks but has remained largely unexplored. By adapting OpenSSL's provider infrastructure to accommodate QKD's pre-distributed key model, we maintain compatibility with current TLS implementations while offering dual layers of security. Performance evaluations demonstrate the feasibility of our hybrid scheme with acceptable overhead, showing that robust security against quantum threats is achievable while addressing the unique requirements of different QKD API specifications.

Index Terms—Post-Quantum Cryptography, POC, OKD, TLS,

[18] https://doi.org/10.48550/arXiv.2503.07196

Recent standardization efforts related to TLS 1.3 have established a framework for hybrid key exchange that combines traditional and post-quantum cryptography [5]. Their approach demonstrates how to achieve security through the concatenation of shared secrets from different key exchange methods, maintaining protection as long as at least one component remains unbroken. This design principle is relevant for QKD deployments where extra security is needed: while QKD provides information-theoretic security, supplementing it with post-quantum cryptography can provide additional protection against implementation vulnerabilities or operational compromises in the OKD system.

Our work explores QKD integration into TLS using OpenSSL's provider infrastructure, encapsulating both PQC shared secrets and QKD key identifiers into a single KEM operation. Unlike previous implementations, we specifically address the integration of stateful OKD protocols through

Quantum Key Distribution for Railwayes an overview of the integration of OKD

Hybrid TLS QKD PO

Authentication in the QKD context



- Two core needs
 - Entity authentication
 - Message/data authentication
- Options
 - PKI-based vs pre-shared symmetric approaches; MACs for data authentication.
- Bootstrapping and lifecycle key management are central design choices

ITU-T Y,380X standards



Table 2.1: Overview of ITU-T Y.380X QKD Standards

Standard	Content
Y.3800	Outlines concepts, capabilities, and design considerations for QKD and QKDN 14.
Y.3801	Defines functional requirements across the quantum, key management, control, and management layers [15].
Y.3802	Specifies the functional architecture, reference points, configurations, and base procedures [16].
Y.3803	Defines the key management architecture, interfaces, and security requirements for QKD networks [17].
Y.3804	Specifies functions and procedures for QKDN control and management via Fault, Configuration, Accounting, Performance, and Security (FCAPS) [18].

The ETSI GS QKD 014 v1.1.1



What are SAEs and KMEs?

- SAE (Secure Application Entity): the application-side client (e.g., encryptors, optical switches, security management systems) that requests keys from its local KME over the ETSI QKD-014 REST/HTTPS API; it sits within the same security boundary/site as its KME, has a unique SAE ID, and can act as master/slave when identical keys must be delivered to multiple peers.
- KME (Key Management Entity): the key-management server in a trusted node that interfaces
 to QKD devices (QKDEs), cooperates with other KMEs across the QKD network, and delivers
 keys to SAEs via a web API; it authenticates SAEs, provides the REST/HTTPS key-delivery
 service, and has a unique KME ID.

Design Idea: Key Storage on Train that uses a "Key petrol station" with ETSI GS QKD 014 v1.1.1

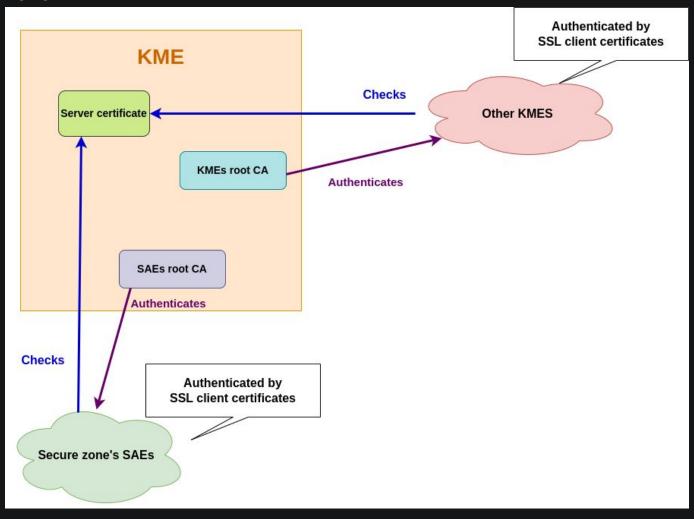




Proof of Concept – Approach Two (3)

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TLS-QKD Topology (2)



TLS variants comparison (3)



35

Scheme / Property	TLS 1.3	TLS-QKD	KEMTLS
Authenticatio n	Signature-based certificate authentication (or Pre-Shared Key). Trust via PKI	possession of a pre shared QKD key proven via a challenge.	Replace signatures with KEMs: certificates contain long-term KEM keys
Hardware Requirements	No specific hardware requirements.	Special QKD system required.	Runs on standard CPUs; no special hardware required.
Resistance to quantum computers?	No, handshake relies on ECDHE for key establishment and ECDSA for authentication.	Yes. Keys have information-theoretic security due to quantum physical effects	Yes, relies only on symmetric primitives.

Design Idea: "Key Petrol Station" ("Schlüsseltankstelle") (4)



ETSI GS QKD 014 v.1.1.1

Phase A: Pre-trip preload

- Stationary QKD setup produces keys in secure location
- Key Management Entity (KME, secure data center with QKD setup) and Secure Application Entity Master (SAE, e.g. the train) authenticate to ensure readiness Get status
- 2. SAE Master requests N slices of keys, KME returns key container Get key
 - a. SAE loads stack of keys into its on-board Hardware Secure Module (HSM)
- KME marks keys so that Slave SAE (e.g. Interlocking) can later call Get key with keyIDs

<u>Phase B — In journey: Use TLS-QKD</u>

<u>Phase C — Refill at stops</u>

Repeat Phase A when docked, replenish and purge expired or used keys.

ETSI GS OKD 014 v1.1.1 https://www.etsi.org/deliver/etsi gs/OKD/001 099/014/01.01.01 60/gs gkd014v010101p.pd